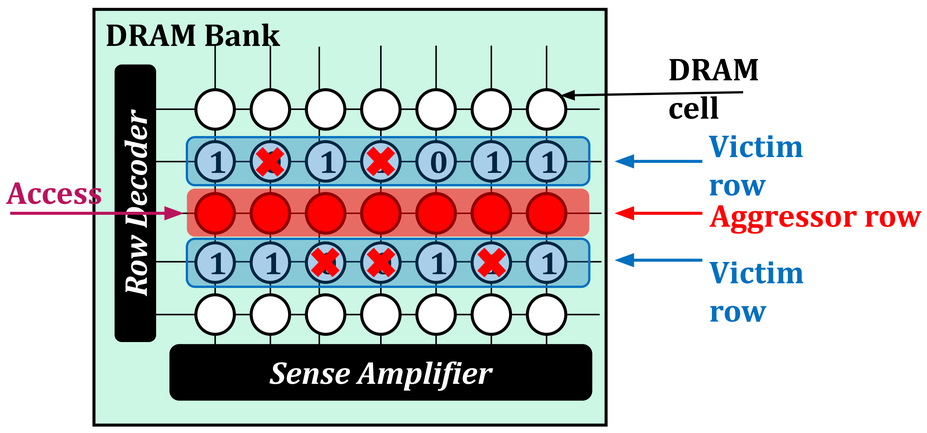
HammerFilter: Robust Protection and Low Hardware Overhead Method for Row-Hammering

Kwangrae Kim, Department of ECE, Hanyang University
Junsu Kim, Department of ECE, Hanyang University
Jeonghyun Woo, Department of ECE, University of Illinois at Urbana-Champaign
Ki-Seok Chung, Department of ECE, Hanyang University

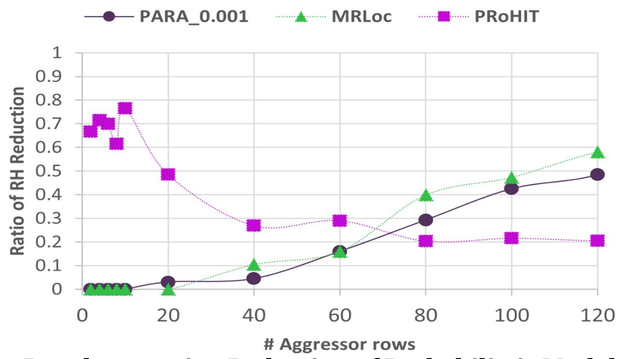
1) Background: Row-hammering



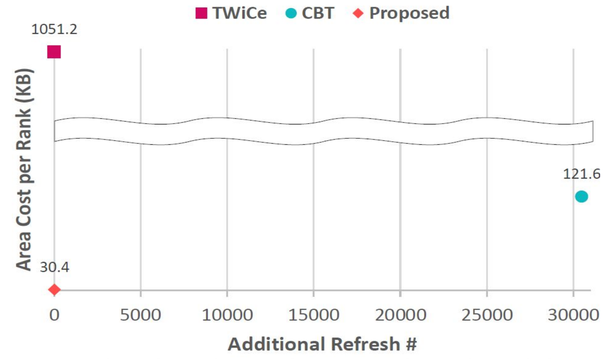
Repeatedly Access to a DRAM row (aggressor row) causes bit-flips in nearby rows (victim row)

2) Motivation

- **The continuous scaling-down of the DRAM process makes DRAM cells more vulnerable to row-hammering**
- **There are two types of hardware-based protection schemes for row-hammering attack: a probabilistic method and a counter-based method**
  1) **The probabilistic schemes have poor protection against complex row-hammering attacks**
     - **PARA[Kim+, ISCA'14], MRLoc[You+, DAC'19], and PRoHIT[Son+, DAC'17] still cannot prevent well in complex row-hammering attacks**
  2) **The counter-based schemes guarantee strong protection, but they suffer from significant area overhead or extreme additional refreshes or even both**
     - **CBT[Seyedzadeh+, ISCA'18] carries out considerable additional refreshes to prevent row-hammering attacks**
     - **TWiCe[Lee+, ISCA'19] incurs significant area overhead to operate it**



Row-hammering Reduction of Probabilistic Models

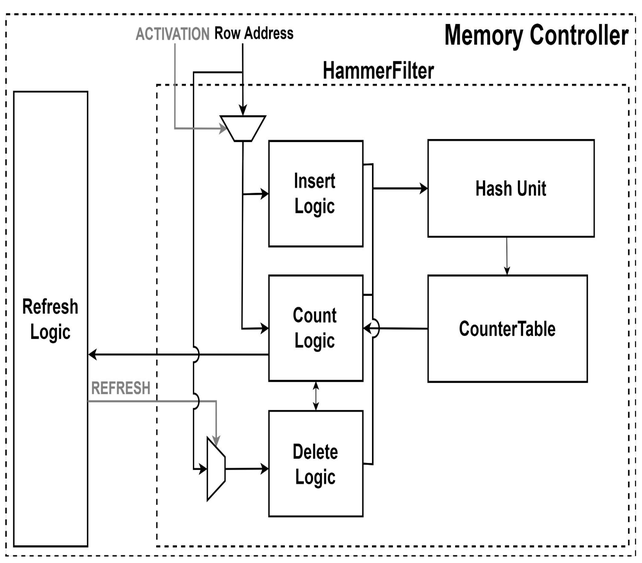


Overhead analysis of Counter-based Models

3) Contribution

1. **We propose a novel data structure, named HammerFilter, which provides guaranteed protection with a low area cost against row-hammering attack**
2. **We evaluated the proposed method with intricate row-hammering attack patterns (five patterns)**
3. **HammerFilter incurs minimal performance overhead in the benign applications**

4) HammerFilter Design Overview

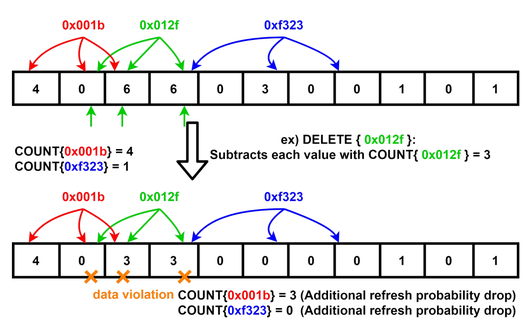


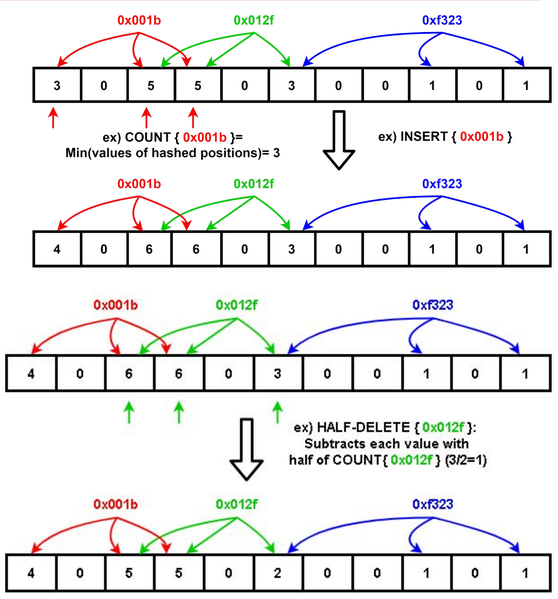
a) Architecture

- **Insert Logic: updates aggressor row candidates with predetermined probability, p_i (0.004)**
- **Delete Logic: reduces corresponding Counter Table's value when additional refresh is done**
- **Count Logic: provides the extent of danger of accessed rows to the Refresh Logic**

b) Description of Operations

- **COUNT: considers the minimum of the values in the hashed positions as the COUNT value**
- **INSERT: increases the values of the hashed positions by one**
- **HALF-DELETE: subtracts each value in the hashed positions by one-half of the COUNT value. It is to avoid a hash-induced collision**



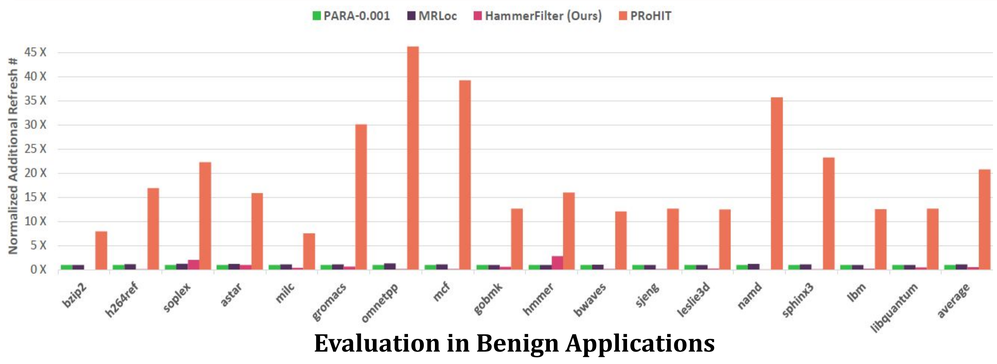


c) HammerFilter's Mechanism

1) **HammerFilter operates INSERT an accessed row to the CounterTable with a fixed probability, p_i.**
2) **It sends COUNT value to the Refresh Logic for every access**
3) **Refresh Logic sends additional refresh with the probability (p_r) according to the COUNT value to the victim rows**

$$\begin{cases} p_r = 1 \div 2^{8-count} & if\ count > 3 \\ p_r = 0 & else \end{cases}$$

5) Experimental Results



Evaluation in Benign Applications

Description of Row-hammering Attack Patterns

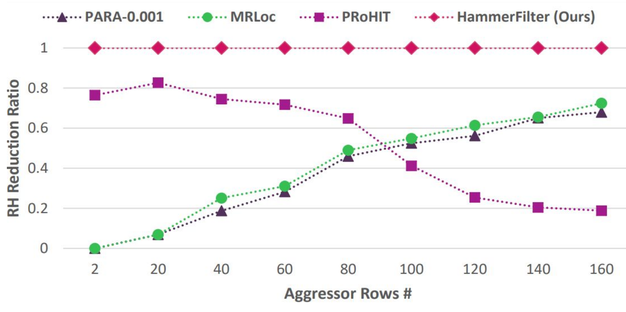
Type	Description	Access Pattern
1	Repeated Selected Rows	$(a_1, a_2, \cdots, a_N)^*$
2	Repeated Selected Rows + Random Rows	$a_1, 1315, a_2, 798, \cdots, a_N, 37, \cdots$
3	Double-sided Rows	$(a_1 - 1, a_1 + 1, \cdots, a_N - 1, a_N + 1)^*$
4	Double-sided Rows + Random Rows	$a_1 - 1, 927, a_1 + 1, \cdots, 109, a_N + 1, \cdots$
5	Double-sided Rows + Repeated Selected Rows	$(a_1 - 1, b_1, a_1 + 1, \cdots, b_N, a_N + 1)^*$

a) Evaluation in Benign Applications

- **The average amount of additional refreshes in HammerFilter is 2.09X and 38.87X less than that of MRLoc and PRoHIT, respectively**

b) Evaluation in Malicious Applications

- **We evaluates HammerFilter with five artificial pattern of row-hammering attack**
- **HammerFilter achieves overwhelmingly better results with respect to all the numbers of aggressor rows than other schemes**



Evaluation in Malicious Applications

6) Conclusion

HammerFilter is a novel method to mitigate row-hammering

- **Prevents all the row-hammering attacks that have various patterns and access multiple rows.**
- **Incurs minimal performance overhead in benign applications**
- **Extremely area efficient compared to counter-based methods**